



HindPhotostat



Hind Photostat & Book Store

Best Quality Classroom Topper Hand Written Notes to Crack GATE, IES, PSU's & Other Government Competitive/ Entrance Exams

MADE EASY

Computer Science Engineering / IT
Toppers Handwritten Notes

Theory Of Computation
By-Ramesh Sundaram Sir

- Theory
- Explanation
- Derivation
- Example
- Shortcuts
- Previous Years Question With Solution

Visit us:-www.hindphotostat.com

Courier Facility All Over India
(DTDC & INDIA POST)

Mob-9311989030



HindPhotostat



MADE EASY , IES MASTER , ACE ACADEMY , KREATRYX

**ESE , GATE, PSU BEST QUALITY TOPPER HAND WRITTEN NOTES
MINIMUM PRICE AVAILABLE @ OUR WEBSITE**

- | | |
|--------------------------------|---------------------------|
| 1. ELECTRONICS ENGINEERING | 2. ELECTRICAL ENGINEERING |
| 3. MECHANICAL ENGINEERING | 4. CIVIL ENGINEERING |
| 5. INSTRUMENTATION ENGINEERING | 6. COMPUTER SCIENCE |

IES , GATE , PSU TEST SERIES AVAILABLE @ OUR WEBSITE

- ❖ IES –PRELIMS & MAINS
- ❖ GATE

➤ **NOTE;- ALL ENGINEERING BRANCHS**

➤ **ALL PSUs PREVIOUS YEAR QUESTION PAPER @ OUR WEBSITE**

PUBLICATIONS BOOKS -

**MADE EASY , IES MASTER , ACE ACADEMY , KREATRYX , GATE ACADEMY , ARIHANT , GK
RAKESH YADAV , KD CAMPUS , FOUNDATION , MC –GRAW HILL (TMH) , PEARSON...OTHERS**

HEAVY DISCOUNTS BOOKS AVAILABLE @ OUR WEBSITE

F230, Lado Sarai New Delhi-110030 Phone: 9311 989 030	Shop No: 46 100 Futa M.G. Rd Near Made Easy Ghitorni, New Delhi-30 Phone:9711475393	F518 Near Kali Maa Mandir Lado Sarai New Delhi-110030 Phone: 9560 163 471	Shop No.7/8 Saidulajab Market Neb Sarai More, Saket, New Delhi-30
--	--	--	--

Website: www.hindPhotostat.com

Contact Us: 9311 989 030

Courier Facility All Over India

(DTDC & INDIA POST)

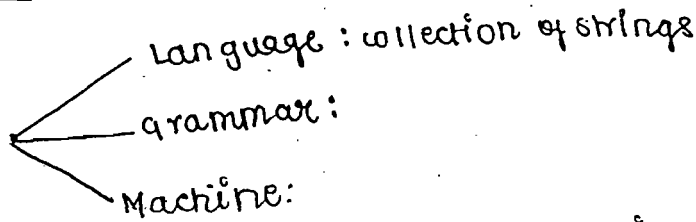
THEORY OF COMPUTATION

- GODEL : Logic is limited
- Turing : Model for computation
- POST
- Chomsky Hierarchy

→ 2 types of computer
Acceptors: Yes/No, given lang. Accept or Not Accept.
Transducer: computational
 x is given $f(x)$ can be computed.

• every problem has associated with language.. we bother about acceptance of language. if we can accept the language.. we can say problem is solvable.

chapter - 0



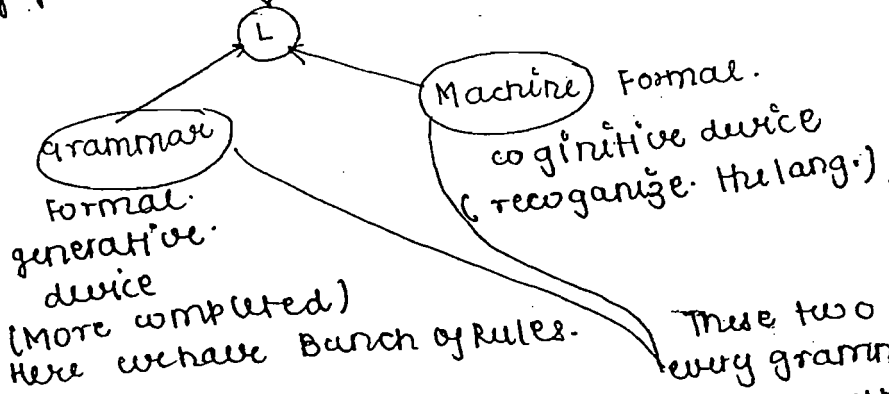
Language can be described by grammar.

- (L) Informal (can't list all the things in language)

(compact) generative.
 its kind of formula.

regular expression

only RE covered by only (RL)



These two generate every grammar.
 (But NOT RE, up to RE)
 because we don't have Machine.

1. Alphabet = $\Sigma = \{a, b\}$
2. string
3. concatenation
4. Reversal
5. length of a string
6. NULL string = "ε"
7. PREFIX
8. SUFFIX

9. substring
10. substring
11. Powers of a string $(w)^n$
12. Σ^* , Σ^+
13. $L \subseteq \Sigma^*$
14. CHOMSKY HEIRARCHY
14. ~~Operatt~~ Representations of Language

• Language

1. alphabet:

14. Representation of Language 6 $\begin{cases} 3 \text{ Formal} \\ 3 \text{ Informal} \end{cases}$

15. operations on language.
union, intersection, L , $L_1 - L_2$, $L_1 \oplus L_2$

16. concatenation of lang.

$$L_1 \cdot L_2 = \{uv \mid u \in L_1, v \in L_2\}$$

17. $L^R = \{u^R \mid u \in L\}$ Reversal of language.

18. L^* , L^+
/ contain ϵ Not contain ϵ

$L^+ = L^* - (\epsilon) \rightarrow$ This is NOT correct statement.

every possible combination of strings.

• Alphabet: a Non-empty finite set of symbols.

- $\Sigma = \{ \}$ Not alphabet
- $\Sigma = \{a\}$, $\Sigma = \{1\}$, $\Sigma = \{2\}$, $\Sigma = \{3\}$ (1 symbol) unary alphabet
- $\Sigma = \{a, b\}$, $\Sigma = \{1, 0\}$ Binary alphabet (2 symbol)
- $\Sigma = \{1, 11, 111, \dots\}$ This is Not allowed, No of symbols should be finite.

$$\{0, 1, 2\} = \{0, 1, 2\}$$

both alphabet same, order dont matter.

symbol $\Sigma = \{ \underline{01}, \underline{10} \}$

compound symbol

$$\Sigma = \{ 01, 10, \textcircled{1}, \textcircled{0} \}$$

This is Not valid symbol
01 or 10 cant break further.

• string: sequence of 0 or more finite symbols taken from the alphabet

sequence: order is important

$$\Sigma = \{a, b\}$$

a^{100} is valid string? \Rightarrow Yes.

$aaa \dots$ 100 times valid.

$aabb$ valid? \Rightarrow No symbols can be taken from alphabet.

$baab = baba$ Not equal string, 'sequence should be follow'

$(ab)^2 \neq a^2b^2$ Not valid in TOC.

$abab \neq aabb$

• concatenation:

if $u = 01$

$v = 100$

$uv = \underline{01} \underline{100}$ concatenation

Here $uv \neq vu$.

for all (u, v)

where $u = 00$

$v = 000$

$uv = vu$ True Here

so Not for all (u, v) ; $uv \neq vu$.

• it is associative

$$u(vw) = (uv)w$$

• Not commutative.

• The length of $u \cdot v$ will always be equal to $(u+v)$.

$u = 01 = 2$

$v = 100 = 3$

$uv = 01100 \Rightarrow$ length is $(2+3) = 5$

• The length of u^k

$u = 100$

$u^k = 001 \Rightarrow u^k \neq u$

for all u and v

$u^k \neq v^k$ False

• $u = u^R$ iff u is a palindrome

Palindrome

- even (because length is integer). : EP
- odd. : OP.

Palindrome language = $\{ \underbrace{ww^R}_{EP} \cup \underbrace{wxw^R}_{OP} \mid w \in \Sigma^*, x \in \Sigma \}$ x is 1 bit

\Downarrow \Downarrow
 1001 10001
 $\{ w \in \Sigma^* \mid w = w^{xc} \}$

Properties of Reversal

$(u^R)^R = u$
 (Reversal of Reversal).
 $(u \cdot v)^R = v^R u^R$
 $(xyz)^R = z^R y^R x^R$

Length:

No of symbols present in strings.

if $\Sigma = \{0, 1, 2\}$.

→ How many length string possible.

0 length → 1 (i.e. ϵ)

1 length → a, b → $|\Sigma| = 2$

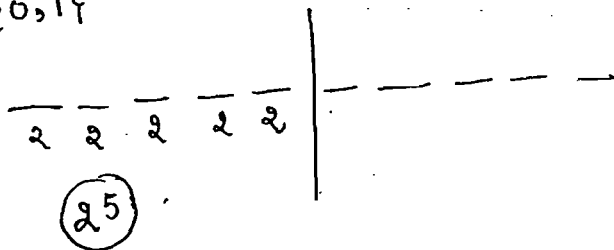
2 length → ab, ba, aa, bb → $|\Sigma|^2$ cardinality of Σ .

⋮

n length → $|\Sigma|^n$.

so of length '3' $2^3 \Rightarrow$ for $\Sigma = \{a, b, c\} \Rightarrow |\Sigma| = 3$
 length '5' 2^5

ww^R How many even palindrome of length 10.
 $\Sigma = \{0, 1\}$



even palindrome = $|\Sigma|^{n/2}$

- odd \Rightarrow length 13. $\Sigma = \{0,1\}$
check even $\rightarrow 12$.

$$|\Sigma| \frac{12}{2} \quad |2|^{12} \Rightarrow 2^6 \times 2 \text{ odd palindrome.}$$

upto length 10, even palindrome. $\Sigma = \{0,1\}$

0	length	$\rightarrow 2^{0/2} = 1$
2	"	$2^{2/2} = 2$
4	"	$2^{4/2} = 4$
8	"	$2^{8/2} = 2^4$
10	"	$2^{10/2}$

• NULL string:

it is only string of length '0'

lang with null string. $|\{\epsilon\}| \rightarrow$ cardinality.
 $|\epsilon| = 0$.
 $|\{\epsilon\}| = 0$.
 empty language.

- language has cardinality.

where
 Null Reverse is Reverse
 Null is identity element for concatenation.

• (Σ^*, \cdot) groupoid
 operator. semi-group.
Algebra.

$$|\Sigma|^5 \rightarrow \Sigma \text{ possible function.}$$

• $|\{0,1\}|^5 \rightarrow |\{0,1\}|$
 how many possible
 all string of length 5.

cardina. = 2^5

$$|B|^{|A|} \text{ function possible} = \underline{\underline{2^{(2^5)}}}$$